Logical Operational Semantics of Parlog
Part II: Or-Parallelism*

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Abstract

This paper refines the definition of a complete, mathematical semantics for the parallel logic programming language PARLOG provided in [3] by giving an explicit Evolving Algebras formalization of the OR-Parallelism in Parlog. In particular we extend the algebras of [3] by new rules which describe the dynamics of the crucial candidate clause search in Parlog which was left abstract in [3].

1 Signature extension of Parlog Algebras

Due to space limitations we skip the Introduction and refer the reader for motivation and basic definition to [1], [2], [3] of the latter of which this paper is a direct sequel. The abstract function candidate-clause(database,node,lit), which in [3] realizes the whole Parlog reduction process of a literal lit yielding as result the body of the selected clause and a unifying substitution of its head and lit, receives here an explicit description by Evolving Algebras rules. To this purpose we extend the previous Parlog Algebras, introducing new universes and functions and modifying the existing ones.

Due to space restrictions we can spell out here only a few new definitions and crucial rules, refering the reader for a full account to [4]. In the full paper [4] we prove the correctness of this candidate-clause specification w.r.t. its definition in [3] by showing that our extended Parlog algebras developed here are "conservative" over the Parlog algebras of [3].

When the Parlog computation system begins the reduction process of a given procedure call lit, it comes into the so-called test-commit-output-spawn phase. In the test phase

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the system tries to find a clause which satisfies the candidate clause condition among those defining the procedure which comes in the form of

\[
\text{procdef(database, lit)} = \underbrace{C_1 \cdot C_{12} \cdots C_{1m}}_{S_1} \cdots \underbrace{C_{n1} \cdot C_{n2} \cdots C_{nm}}_{S_n}
\]

The parclauses \( S_i \) are searched through sequentially (seq-search). Each parclause is a sequence \( C_1.C_2.\cdots.C_m \) of guarded clauses which are searched through in parallel (or-search). The try-clause computation, whether a clause is a candidate clause for \( \text{lit} \), checks whether the clause is an unguarded clause for which input matching\(^1\) of \( \text{lit} \) and head succeeds, or whether it is a guarded clause for which: a) the input-match computation (unification of \( \text{lit} \) and head) succeeds with a substitution \( s \); b) the guard-eval computation of the clause guard by the program succeeds with a substitution \( s' \); c) the substitutions \( s \) and \( s' \) are consistent. (seq-search, or-search, try-clause, input-match, guard-eval are new tags.)

When a candidate clause is found, the calling literal \( \text{lit} \) commits to it (commit phase) - interrupting the search for (other) candidate clauses -; output unification is performed between the output mode argument terms of \( \text{lit} \) and those of the selected clause (output phase); the \( \text{lit} \) computation is reduced to the evaluation of the candidate clause's body under the computed output substitution (spawn phase).

For an explicit description of this candidate clause search, we imagine each or-par node as root of a computation subtree performing the test-commit-output phase and of another (later created) subtree to perform the spawn phase.

During the input-match computation a node may come into a new mode suspended, namely when there is a unifier for the calling literal and the clause head which however tries to bind a variable occurring (in a term) in an input argument position of the call. During the spawn phase we need a function \( \text{sub} : \text{Node} \rightarrow \text{Sub} \) which yields a substitution which is known at the given node, but which is not transparent to the main system. In the or-par node subtree description, the function goal will be used to pass from parent to child the calling literal \( \text{lit} \) which is responsible for the candidate clause search. The same function will be used to report from child to parent the body of the identified candidate clause for the calling literal.

Some other new mostly self explanatory functions will be introduced where needed in the rules.

2 The or-par node operation

The transition rules of Part I (see [3]) that perform the or-par node operation, are replaced by those that we are describing in this section.

When the or-par node has mode starting, it has just been created as child of a node

\(^1\)Input matching (of a literal \( \text{lit} \) and a term \( t \)) is defined as unification of \( \text{lit} \) and \( t \) in which no variable is substituted which occurs in an input argument of \( \text{lit} \) (with respect to that program).

In contrast to input matching one speaks of output unification to refer to a unification of two terms which appear in an output mode argument (of a literal w.r.t. a given program). This façon de parler stresses that for the unification of terms occurring in output mode arguments there is no restriction on the direction of the bindings (from goal to clause head (output) or viceversa (input)).