
Chapter 31

Non-cooperative game theory

31.1 In an n -person game we assign to each player i ($i = 1, \dots, n$) a *strategy set* S_i and a pure strategy payoff function u_i that gives player i utility $u_i(\mathbf{s}) = u_i(s_1, \dots, s_n)$ for each *strategy profile* $\mathbf{s} = (s_1, \dots, s_n) \in S = S_1 \times \dots \times S_n$.

An n -person game in *strategic* (or *normal*) form. If all the strategy sets S_i have a finite number of elements, the game is called *finite*.

31.2 A strategy profile (s_1^*, \dots, s_n^*) for an n -person game is a *pure strategy Nash equilibrium* if for all $i = 1, \dots, n$ and all s_i in S_i ,
$$u_i(s_1^*, \dots, s_n^*) \geq u_i(s_i, s_1^*, \dots, s_{i-1}^*, s_{i+1}^*, \dots, s_n^*)$$

Definition of a pure strategy Nash equilibrium for an n -person game.

31.3 If for all $i = 1, \dots, n$, the strategy set S_i is a nonempty, compact, and convex subset of \mathbb{R}^m , and $u_i(s_1, \dots, s_n)$ is continuous in $S = S_1 \times \dots \times S_n$ and quasiconcave in its i th variable, then the game has a pure strategy Nash equilibrium.

Sufficient conditions for the existence of a pure strategy Nash equilibrium. (There will usually be several Nash equilibria.)

31.4 Consider a finite n -person game where S_i is player i 's pure strategy set, and let $S = S_1 \times \dots \times S_n$. Let Ω_i be a set of probability distributions over S_i . An element σ_i of Ω_i (σ_i is then a function $\sigma_i : S_i \rightarrow [0, 1]$) is called a *mixed strategy* for player i , with the interpretation that if i plays σ_i , then i chooses the pure strategy s_i with probability $\sigma_i(s_i)$.

Definition of a *mixed strategy* for an n -person game.

If the players choose the *mixed strategy profile* $\boldsymbol{\sigma} = (\sigma_1, \dots, \sigma_n) \in \Omega_1 \times \dots \times \Omega_n$, the probability that the pure strategy profile $\mathbf{s} = (s_1, \dots, s_n)$ occurs is $\sigma_1(s_1) \dots \sigma_n(s_n)$. The expected payoff to player i is then

$$u_i(\boldsymbol{\sigma}) = \sum_{\mathbf{s} \in S} \sigma_1(s_1) \dots \sigma_n(s_n) u_i(\mathbf{s})$$

- 31.5 A mixed strategy profile $\sigma^* = (\sigma_1^*, \dots, \sigma_n^*)$ is a *Nash equilibrium* if for all i and every σ_i ,
- $$u_i(\sigma^*) \geq u_i(\sigma_1^*, \dots, \sigma_{i-1}^*, \sigma_i, \sigma_{i+1}^*, \dots, \sigma_n^*)$$
- Definition of a *mixed strategy Nash equilibrium* for an n -person game.
- 31.6 σ^* is a Nash equilibrium if and only if the following conditions hold for all $i = 1, \dots, n$:
- $$\begin{aligned} \sigma_i^*(s_i) > 0 &\Rightarrow u_i(\sigma^*) = u_i(s_i, \sigma_{-i}^*) \quad \text{for all } s_i \\ \sigma_i^*(s'_i) = 0 &\Rightarrow u_i(\sigma^*) \geq u_i(s'_i, \sigma_{-i}^*) \quad \text{for all } s'_i \end{aligned}$$
- where $\sigma_{-i}^* = (\sigma_1^*, \dots, \sigma_{i-1}^*, \sigma_{i+1}^*, \dots, \sigma_n^*)$ and we consider s_i and s'_i as degenerate mixed strategies.
- An equivalent definition of a (mixed strategy) Nash equilibrium.
- 31.7 Every finite n -person game has a mixed strategy Nash equilibrium.
- An important result.
- 31.8 The pure strategy $s_i \in S_i$ of player i is *strictly dominated* if there exists a mixed strategy σ_i for player i such that for all feasible combinations of the other players' pure strategies, i 's payoff from playing strategy s_i is strictly less than i 's payoff from playing σ_i :
- $$\begin{aligned} u_i(s_1, \dots, s_{i-1}, s_i, s_{i+1}, \dots, s_n) < \\ u_i(s_1, \dots, s_{i-1}, \sigma_i, s_{i+1}, \dots, s_n) \end{aligned}$$
- for every $(s_1, \dots, s_{i-1}, s_{i+1}, \dots, s_n)$ that can be constructed from the other players' strategy sets $S_1, \dots, S_{i-1}, S_{i+1}, \dots, S_n$.
- Definition of strictly dominated strategies.
- In an n -person game, the following results hold:
- 31.9
- If iterated elimination of strictly dominated strategies eliminates all but the strategies (s_1^*, \dots, s_n^*) , then these strategies are the unique Nash equilibrium of the game.
 - If the mixed strategy profile $(\sigma_1^*, \dots, \sigma_n^*)$ is a Nash equilibrium and, for some player i , $\sigma_i^*(s_i) > 0$, then s_i survives iterated elimination of strictly dominated strategies.
- Useful results.
Iterated elimination of strictly dominated strategies need not result in the elimination of any strategy. (For a discussion of iterated elimination of strictly dominated strategies, see the literature.)