

# Singularity Theory for Nonlinear Optimization Problems

J. Casti

*International Institute for Applied Systems Analysis, A-2361 Laxenburg, Austria*

## 1. Background

Consider a smooth ( $C^\infty$ ) function  $f: R^n \rightarrow R^m$  and assume that  $f$  has a critical point at the origin, i.e.,  $df(0) = 0$ . The theory of singularities as developed by Thom, Mather, Arnol'd, and others (Lu, 1976; Gibson, 1979; Arnol'd, 1981) addresses the following basic questions:

- (1) What is the local character of  $f$  in a neighborhood of the critical point? Basically, this question amounts to asking "at what point is it safe to truncate the Taylor series for  $f$ ?" This is the *determinacy* problem.
- (2) What are the "essential" perturbations of  $f$ ? That is, what perturbations of  $f$  can occur that change the qualitative nature of  $f$  and that cannot be transformed away by a change of coordinates? This is the *unfolding* problem.
- (3) Can we classify the types of singularities that  $f$  can have up to diffeomorphism? This is the *classification* problem.

Elementary catastrophe theory largely solves these three problems (when  $m = 1$ ); its generalization to singularity theory solves the first two, and gives relatively complete information on the third for  $m, n$  small. Here we outline a program for the utilization of these results in an applied setting to deal with certain types of nonlinear optimization problems. In the following section we give a brief summary of the main results of singularity theory for problems (1)–(3) for *functions* ( $m = 1$ ) and then proceed to a discussion of how these results may be employed for nonlinear optimization.

## 2. Determinacy, Unfoldings, and Classifications

### 2.1. Equivalence of germs

In its local version, elementary catastrophe theory deals with functions  $f: U \rightarrow R$  where  $U$  is a neighborhood of  $O$  in  $R^n$ . The cleanest way to handle

such functions is to pass to *germs*, a germ being a class of functions that agrees on suitable neighborhoods of  $O$ . All operations on germs are defined by performing similar operations on representatives of their classes. In the sequel, we usually make no distinction between a germ and a representative function.

We let  $E_n$  be the set of all smooth germs  $R^n \rightarrow R$ , and let  $E_{nm}$  be the set of all smooth germs  $R^n \rightarrow R^m$ . Of course  $E_{n,1} = E_n$ . These sets are vector spaces over  $R$ , of infinite dimension. We abbreviate  $(x_1, \dots, x_n) \in R^n$  to  $x$ . If  $f \in E_{nm}$  then

$$f(x) = [f_1(x), \dots, f_m(x)]$$

and the  $f_i$  are the *components* of  $f$ .

A *diffeomorphism germ*  $\varphi: R^n \rightarrow R^n$  satisfies  $\varphi(0) = 0$ , and has an inverse  $\varphi'$  such that  $\varphi[\varphi'(x)] = x = \varphi'[\varphi(x)]$  for  $x$  near 0. It represents a smooth, invertible local coordinate change. By the Inverse Function Theorem,  $\varphi$  is a diffeomorphism germ if and only if it has a nonzero Jacobian, that is,

$$\det[\partial\varphi_i / \partial x_j(0)] \neq 0$$

Two germs  $f, g: R^n \rightarrow R$  are *right equivalent* if there is a diffeomorphism germ  $\varphi$  and a constant  $\gamma \in R$  such that

$$g(x) = f(\varphi(x)) + \gamma$$

This is the natural equivalence for studying topological properties of the gradient  $\nabla f$  (Poston and Stewart, 1978). If  $f$ , rather than  $\nabla f$ , is important, the term  $\gamma$  is omitted.

A *type* of germ is a right equivalence class and the classification of germs up to right equivalence amounts to a classification of types. Each type forms a subset of  $E_n$ , and the central object of study is the way these types fit together.

A precise description is complicated by the fact that most types have infinite dimension; but there is a measure of the complexity of a type, the *codimension*, which is generally finite. Heuristically, it is the difference between the dimension of the type and that of  $E_n$  (even though both are infinite). A precise definition is given below.

The largest types have codimension 0 and form open sets in  $E_n$ . Their boundaries contain types of codimension 1; the boundaries of these in turn contain types of codimension 2, and so on, with higher codimensions revealing progressively more complex types. Types of infinite codimension exist, but form a very small set in a reasonable sense.